
Graphs and Algorithms

Exercise 1 (Vertex vs. edge connectivity)

For all $t \in \mathbb{N}$ show that there exists a graph with vertex connectivity 1 and edge connectivity t .

Exercise 2 (Exam question, 2010.)

Let $k \geq 1$ and let G be a graph on n vertices such that any two distinct vertices x and y satisfy $\deg(x) + \deg(y) \geq n + k - 2$. Prove that G is k -connected.

Exercise 3 (The k -dimensional grid)

Is the k -dimensional grid (as defined in exercise 5 of sheet 2), k -connected?

Exercise 4 (Dilworth's theorem)

Let (P, \leq) , where P is a finite set, be a partially ordered set. To recapitulate, (P, \leq) is a partially ordered set (or *poset*) if the relation \leq on P satisfies

1. $a \leq a$,
2. $a \leq b$ and $b \leq a$ implies $a = b$ and
3. $a \leq b$ and $b \leq c$ implies $a \leq c$.

Note that we do *not* require that for all pairs (a, b) one of $a \leq b$ and $b \leq a$ holds, as we would do for a total ordering.

We say that a subset $C \subseteq P$ is a *chain* if $a \leq b$ or $b \leq a$ for every two elements $a, b \in C$. Similarly, we call a subset $A \subseteq P$ an *antichain* if neither $a \leq b$ nor $b \leq a$ for every two distinct elements $a, b \in A$. Finally, we say that a set C_1, \dots, C_t of chains covers P if every element of P belongs to some chain C_i .

Prove that the maximum number of elements in an antichain equals to the minimum number of chains needed to cover P .

Hint 1: Prove that we need at least as many chains as there are elements in an antichain. (This is easy.)

Hint 2: Consider the auxiliary bipartite graph $G = (A \cup B, E)$, where A and B are disjoint copies of P , with an edge between $a \in A$ and $b \in B$ if $a \leq b$ and a, b represent distinct elements in P . Construct an antichain in (P, \leq) from each vertex cover in G . What about matchings in G and sets of covering chains?

DISCUSSION OF THE SOLUTION IN THE EXERCISE CLASS ON 21.03.2013.