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## Algorithmic Game Theory HS 2013

### Exercise sheet 10

#### EXERCISE 10.1:

In this exercise we consider combinatorial auctions with single-minded bidders. Recall that in such an auction every player is only interested in getting the goods in  $S_i \subseteq U$  (where  $U$  is the set of goods). The player  $i$  values this bundle  $S_i$  with  $v_i \in \mathbb{R}^+$ . Both  $S_i$  and  $v_i$  are the private information of player  $i$ . Every player  $i$  submits a bid  $(B_i, b_i)$  to the auction, expressing the desire to get the bundle  $B_i$  and that the player values it with  $b_i$ .

Recall the characteristics of VCG and the LOS mechanisms. In VCG the mechanism computes an optimal allocation  $\{S_i^*\}_{i=1}^n$  of goods to the players (where the allocation maximizes the sum of the valuations of all players), and the payment  $p_i$  to every player  $i$ :

$$p_i = \sum_{j \neq i} b_j(\bar{S}_j) - \sum_{j \neq i} b_j(S_j^*),$$

where  $\{\bar{S}_j\}_{j \neq i}$  is an assignment maximizing the total valuation of players  $1, 2, \dots, i-1, i+1, \dots, n$ .

In a LOS mechanism a greedy algorithm is used to compute an approximate solution. In each iteration it grants the bid with the highest value according to the formula  $b_i/\sqrt{|B_i|}$ , after which it removes the bids that are blocked by  $B_i$  before reiterating. The payment to a player  $i$  is then  $q_i = b_j \sqrt{|B_i|/|B_j|}$ , where player  $j$  is the highest uniquely blocked bidder of  $i$ . In both mechanisms a player who is not granted his bid pays nothing.

- a) Consider the VCG mechanism and the LOS mechanism for a combinatorial auction with single-minded bidders.

Provide a problem instance for each one of the following settings:

- i) The total sum of payments in the VCG mechanism is greater than the total sum of payments in the LOS mechanism.
  - ii) The total sum of payments in the LOS mechanism is greater than the total sum of payments in the VCG mechanism.
- b) Consider the LOS greedy algorithm for granting bids of players. In the lecture we have seen that a player  $i$  with her bid  $(B_i, b_i)$  can uniquely block (u-block for short) a player  $j$  with her bid  $(B_j, b_j)$  even if  $B_i \cap B_j = \emptyset$ . Show, however, that if  $j$  is the highest u-blocked bid by player  $i$ , then  $B_i \cap B_j \neq \emptyset$ .
- c) Consider the following modification of the LOS mechanism:
- i) the outcome (i.e., the decision of the mechanism about which player is granted its bundle) remains unchanged;
  - ii) the price that any winner  $i$  pays is  $\sqrt{|B_i|} \frac{b_j}{\sqrt{|B_j|}}$ , where  $j > i$  is the first  $j$  after  $i$  (in the order given by the descending values of  $b_k/\sqrt{|B_k|}$ ,  $k = 1, \dots, n$ ) for which  $B_i \cap B_j \neq \emptyset$ . The payment will be zero if no such  $j$  exists.

Is this mechanism truthful?